

another proof of cayley's formula

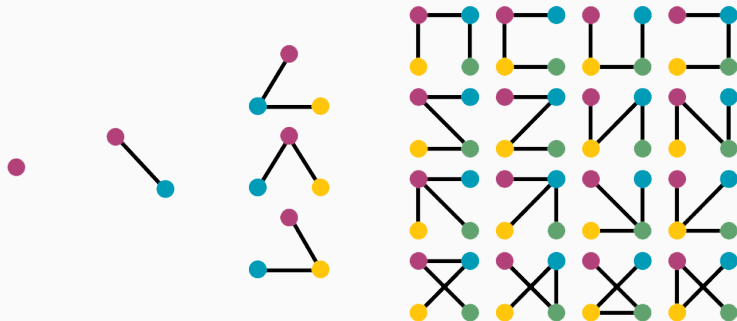
an introduction to analytic combinatorics

Eva Hainzl, 2019

Cayley's formula

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For every $n \in \mathbb{N}$, the number of labelled trees on n vertices is n^{n-2} .



Cayley's formula

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For every $n \in \mathbb{N}$, the number of labelled trees on n vertices is n^{n-2} .

Equivalent problem

For every $n \in \mathbb{N}$, the number of labelled, rooted trees on n vertices is n^{n-1} .



1. Find the generating function (implicit expression)
Symbolic method
2. Analysis of the generating function
Lagrange inversion theorem

Definition

A combinatorial class, or simply a class, is a finite or denumerable set on which a size function is defined, satisfying the following conditions:

1. the size of an element is a non-negative integer
2. the number of elements of any given size is finite

Definition

The counting sequence of a class \mathcal{A} is the sequence of integers $(A_n)_{n \geq 0}$ where A_n is the number of objects in class \mathcal{A} that have size n .

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The ordinary generating function (OGF) of a combinatorial class \mathcal{A} is the formal power series

$$A(z) = \sum_{n=0}^{\infty} A_n z^n$$

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Definition

The exponential generating function (EGF) of a combinatorial class \mathcal{A} is the formal power series

$$A(z) = \sum_{n=0}^{\infty} \frac{A_n}{n!} z^n \quad A_n = n! \cdot [z^n]A(z)$$

The Symbolic Method

The Symbolic Method

Example (Binary Trees (rooted, plane))

$n=1$



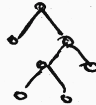
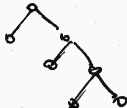
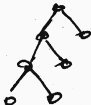
$n=3$



$n=5$



$n=7$



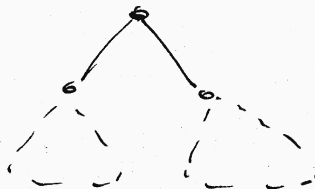
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The Symbolic Method

Example (Binary Trees (rooted, plane))

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or



The Symbolic Method

Example (Binary Trees (rooted, plane))

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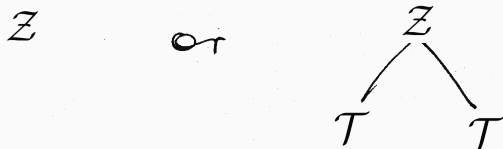
or



- Class of binary trees, $\mathcal{T} = \{\text{binary trees}\}$, size = # of vertices

The Symbolic Method

Example (Binary Trees (rooted, plane))



- Class of binary trees: $\mathcal{T} = \{\text{binary trees}\}$, size = # of vertices
- Atomic class: $\mathcal{Z} = \{\bullet\}$, size = 1

The Symbolic Method

Example (Binary Trees (rooted, plane))



- Class of binary trees: $\mathcal{T} = \{\text{binary trees}\}$, size = # of vertices
- Atomic class: $\mathcal{Z} = \{\bullet\}$, size = 1
- Neutral class: $\mathcal{E} = \{\text{empty graph}\}$, size = 0

The Symbolic Method

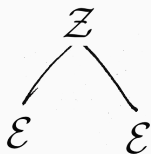
Example (Binary Trees (rooted, plane))



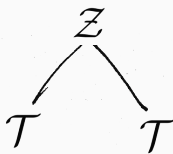
$$\mathcal{T} = \mathcal{Z} \times \mathcal{E} \times \mathcal{E} + \mathcal{Z} \times \mathcal{T} \times \mathcal{T}$$

The Symbolic Method

Example (Binary Trees (rooted, plane))



or



$$\mathcal{T} = \mathcal{Z} \times \mathcal{E} \times \mathcal{E} + \mathcal{Z} \times \mathcal{T} \times \mathcal{T}$$

$$T(z) = z \cdot 1 \cdot 1 + z \cdot T(z) \cdot T(z) = z(1 + T(z)^2)$$

Observation

1. The EGF of \mathcal{E} is 1.
 \mathcal{E} contains by definition only a single object of size 0.

The Symbolic Method

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The Symbolic Method

Observation

1. The EGF of \mathcal{E} is 1.
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2. The EGF of \mathcal{Z} is z .
 \mathcal{Z} contains by definition only a single object of size 1.
3. The EGF of $\mathcal{A} + \mathcal{B}$ is $A(z) + B(z)$, given $\mathcal{A} \cap \mathcal{B} = \emptyset$.
If $\mathcal{A} \cap \mathcal{B} = \emptyset$ then we have either an element from \mathcal{A} or from \mathcal{B} .
Count elements of size n in \mathcal{A} then in \mathcal{B} and add them up.

The Symbolic Method

Observation

The EGF of $\mathcal{C} = \mathcal{A} \times \mathcal{B}$ is $A(z) \cdot B(z)$.

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Let $c \in \mathcal{A} \times \mathcal{B}$, $\text{size}(c) = n$.

Then c is a combination of a and b ,

where $a \in \mathcal{A}$, $\text{size}(a) = k$ and $b \in \mathcal{B}$, $\text{size}(b) = n - k$.

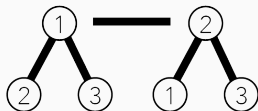
\Rightarrow True for OGF's.

The Symbolic Method

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The EGF of $\mathcal{C} = \mathcal{A} \times \mathcal{B}$ is $A(z) \cdot B(z)$.

When we combine objects, we have to relabel them...

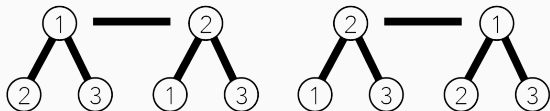


The Symbolic Method

Observation

The EGF of $\mathcal{C} = \mathcal{A} \times \mathcal{B}$ is $A(z) \cdot B(z)$.

But we can't just relabel randomly, otherwise we doublecount!

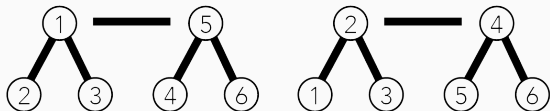


The Symbolic Method

Observation

The EGF of $\mathcal{C} = \mathcal{A} \times \mathcal{B}$ is $A(z) \cdot B(z)$.

Relabeling has to be order-preserving

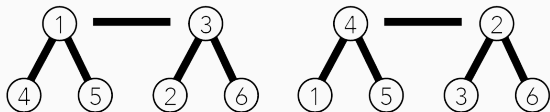


The Symbolic Method

Observation

The EGF of $\mathcal{C} = \mathcal{A} \times \mathcal{B}$ is $A(z) \cdot B(z)$.

Choose k labels for left object, use the other $(n-k)$ labels for right object



$$C_n = \sum_{k=0}^n \binom{n}{k} A_k B_{n-k} = n! \sum_{k=0}^n \frac{A_k}{k!} \cdot \frac{B_{n-k}}{(n-k)!}$$

The Symbolic Method

Sequence Construction

-

$$\mathcal{A} = \mathcal{Z}$$

The Symbolic Method

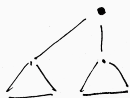
Sequence Construction



$$\mathcal{A} = \mathcal{Z} \times (\mathcal{E} + \mathcal{A})$$

The Symbolic Method

Sequence Construction



$$\mathcal{A} = \mathcal{Z} \times (\mathcal{E} + \mathcal{A} + \mathcal{A} \times \mathcal{A})$$

The Symbolic Method

Sequence Construction



$$\mathcal{A} = \mathcal{Z} \times (\mathcal{E} + \mathcal{A} + \mathcal{A} \times \mathcal{A} + \mathcal{A} \times \mathcal{A} \times \mathcal{A})$$

The Symbolic Method

Sequence Construction



$$\mathcal{A} = \mathcal{Z} \times (\mathcal{E} + \mathcal{A} + \mathcal{A} \times \mathcal{A} + \mathcal{A} \times \mathcal{A} \times \mathcal{A} + \dots)$$

The Symbolic Method

Sequence Construction



$$\begin{aligned}\mathcal{A} &= \mathcal{Z} \times (\mathcal{E} + \mathcal{A} + \mathcal{A} \times \mathcal{A} + \mathcal{A} \times \mathcal{A} \times \mathcal{A} + \dots) \\ &= \mathcal{Z} \times \text{SEQ}(\mathcal{A})\end{aligned}$$

The Symbolic Method

Sequence Construction



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$$A(z) = z \cdot \sum_{n \geq 0} A(z)^n = z \cdot \frac{1}{1 - A(z)}$$

The Symbolic Method

Set Construction



← UNORDERED!

$$\mathcal{A} = \mathbb{Z} \times \text{SET}(\mathcal{A})$$

The Symbolic Method

Set Construction



$$\mathcal{A} = \mathcal{Z} \times \text{SET}(\mathcal{A})$$

$$A(z) = z \cdot \left(1 + A(z) + \frac{A(z)^2}{2} + \frac{A(z)^3}{3!} + \frac{A(z)^4}{4!} \dots \right) = z \cdot \exp(A(z))$$

Summary

If we can (literally) describe a combinatorial class in terms of "and, or, set, sequence", we can immediately derive the generating function!

Construction	Symbolic	EGF
Sum	$\mathcal{A} + \mathcal{B}$	$A(z) + B(z)$
Product	$\mathcal{A} \times \mathcal{B}$	$A(z) \cdot B(z)$
Sequence	$\text{SEQ}(\mathcal{A})$	$(1 - A(z))^{-1}$
Set	$\text{SET}(\mathcal{A})$	$\exp(A(z))$

Analysis of the EGF

Singularity Analysis

Theorem (Cauchy-Hadamard)

Let $A(z) = \sum_{n \geq 1} A_n(z - c)^n$ be a power series with $A_n, c \in \mathbb{C}$. Then the radius of convergence of A at the point c is given by

$$\frac{1}{R} = \limsup_{n \rightarrow \infty} |A_n|^{\frac{1}{n}}$$

Singularity Analysis

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$$\frac{1}{R} = \limsup_{n \rightarrow \infty} |A_n|^{\frac{1}{n}}$$

$$|A_n| \sim R^{-n} \phi(n)$$

\implies "Singularity analysis"

Cauchy Integral Formula

Theorem (Cauchy Integral Formula)

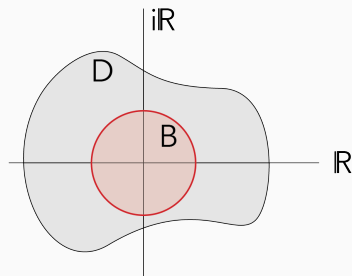
Let f be a holomorph function in an open subset $D \subseteq \mathbb{C}$ and let $\overline{B_r(c)} \subseteq D$. Then it holds for all $z \in B_r(c)$:

$$f(z) = \frac{1}{2\pi i} \int_{\partial B} \frac{f(y)}{y - z} dy$$

As a consequence, f is analytic in $B_r(c)$ and for its Taylor series $f(z) = \sum f_n(z - c)^n$ it holds that

$$f_n = \frac{1}{2\pi i} \int_{\partial B} \frac{f(y)}{(y - c)^{n+1}} dy$$

Cauchy Integral Formula



If f is "nice" in D then it has Taylor series $f(z) = \sum f_n(z - c)^n$ in B and

$$f_n = \frac{1}{2\pi i} \int_{\partial B} \frac{f(y)}{(y - c)^{n+1}} dy$$

\implies Basis for most basic theorems in analytic combinatorics

Lagrange Inversion Theorem

Theorem (Lagrange Inversion Theorem)

Let $\phi(u) = \sum_{k \geq 0} \phi_k u^k$ be a power series of $\mathbb{C}[[u]]$ with $\phi_0 \neq 0$. Then, the equation $A = z\phi(A)$ admits a unique solution in $\mathbb{C}[[u]]$ whose coefficients are given by

$$A(z) = \sum_{n \geq 1} A_n z^n, \quad \text{where } A_n = \frac{1}{n} [u^{n-1}] \phi(u)^n$$

Cayley's formula



$$\mathcal{A} = \mathcal{Z} \times \text{SET}(\mathcal{A}) \implies A(z) = \sum_{n \geq 0} \frac{A_n}{n!} z^n = z \cdot \exp(A(z))$$

Cayley's formula



$$\mathcal{A} = \mathcal{Z} \times \text{SET}(\mathcal{A}) \implies A(z) = \sum_{n \geq 0} \frac{A_n}{n!} z^n = z \cdot \exp(A(z))$$

By Lagrange Inversion Theorem:

$$\begin{aligned} \frac{A_n}{n!} &= [z^n] A(z) = \frac{1}{n} [u^{n-1}] \exp(u)^n \\ &= \frac{1}{n} [u^{n-1}] \sum_{n \geq 0} \frac{(nu)^n}{n!} = \frac{n^{n-1}}{n(n-1)!} = \frac{n^{n-1}}{n!} \end{aligned}$$

Find out more:

Philippe Flajolet and Robert Sedgewick, *Analytic Combinatorics*,
Cambridge University Press (2009).

Available online: <http://algo.inria.fr/flajolet/Publications/book.pdf>