# The interpretability lattice of clonoids is distributive 

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## Clonoids (AKA minions AKA minor closed sets)

- A functional clonoid $\mathcal{C}$ on sets $A, B$ is a nonempty family of operations from $A$ to $B$ closed under taking minors
- Taking minors: $\sigma:[n] \rightarrow[m]$ sends $n$-ary $f$ to $m$-ary $f^{\sigma}$ where $f^{\sigma}\left(x_{1}, \ldots, x_{m}\right)=f\left(x_{\sigma(1)}, \ldots, x_{\sigma(n)}\right)$


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## Clonoid homomorphisms

- $\phi: \mathcal{C} \rightarrow \mathcal{D}$ preserves arity and commutes with taking minors
- Another view: homomorphism sends identities true in $\mathcal{C}$ to identities of $\mathcal{D}$ - can interpret $\mathcal{C}$ in $\mathcal{D}$
- Example:

$$
f(x, x, y) \approx g(x, y, y, z) \Rightarrow \phi(f)(x, x, y) \approx \phi(g)(x, y, y, z)
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- For $\mathbb{A}, \mathbb{B}, \mathbb{A}^{\prime}, \mathbb{B}^{\prime}$ finite relational structures $\operatorname{Pol}(\mathbb{A}, \mathbb{B}) \rightarrow \operatorname{Pol}\left(\mathbb{A}^{\prime}, \mathbb{B}^{\prime}\right)$ gives a reduction from $\operatorname{PCSP}\left(\mathbb{A}^{\prime}, \mathbb{B}^{\prime}\right)$ to $\operatorname{PCSP}(\mathbb{A}, \mathbb{B})$ [Bulín, Krokhin, Opršal; 2018]


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## Interpretability lattice

- $\mathcal{C} \rightarrow \mathcal{D}$ is a quasiorder - factorize and look at the poset $\mathcal{L}$
- Libor Barto: A lot of categories embed (fully) into $\mathcal{L}$ (it is alg-universal)
- Warning: $\mathcal{L}$ is class-size
- $\mathcal{L}$ restricted to finite clonoids: continuum-sized
- $\mathcal{L}$ is also a lattice


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## Meet

- $\mathcal{C}$ goes from $A_{1}$ to $B_{1}$ and $\mathcal{D}$ goes from $A_{2}$ to $B_{2}$
- $\mathcal{C} \wedge \mathcal{D}$ has operations $(f, g)$ where $f$ and $g$ are operations in $\mathcal{C}$ and $\mathcal{D}$ - $(f, g): A_{1}^{n} \times A_{2}^{n} \rightarrow B_{1} \times B_{2}$ is defined componentwise


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## Join

- Abstract clonoids: $\mathcal{C} \vee \mathcal{D}$ satisfies all identities of $\mathcal{C}$ and $\mathcal{D}$
- Operations of $\mathcal{C} \vee \mathcal{D}$ : Disjoint union of $\mathcal{C}$ and $\mathcal{D}$
- $\mathcal{C} \vee \mathcal{D}$ goes from $A_{1} \cup A_{2}$ to $B_{1} \cup B_{2} \cup\{\star\}$, each operation comes either from $\mathcal{C}$ or from $\mathcal{D}$.
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## Distributivity

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\mathcal{C} \wedge(\mathcal{D} \vee \mathcal{E})=(\mathcal{C} \wedge \mathcal{D}) \vee(\mathcal{C} \wedge \mathcal{E})
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## Distributive sublattices of $\mathcal{L}$

- Any finite distributive lattice is a sublattice of $\mathcal{L}$
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Thank you for your attention.

