Deciding the existence of a *k*-wnu operation in polynomial time

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Maltsev conditions

• A strong Maltsev condition is a system of equalities for some operations and variables, eg.

 $p(p(x, y), r(y)) \approx x.$

- A variety V satisfies the Maltsev condition M if we can replace the operations in M by terms of V and get a system of equalities true in V.
- Linear Maltsev: No nested operations. Example:

 $m(x, x, y) \approx y$ $m(y, x, x) \approx y.$

• Libor Barto: CSP complexity can be characterized by linear Maltsev conditions.

- Given finite **A**, we want to know if V(**A**) satisfies a given fixed Maltsev condition.
- Example: "Does A have a Maltsev term?" (Complexity: Unknown.)
- Known to be EXPTIME-complete for numerous common Maltsev conditions (Freese, Valeriote, Horowitz).
- Restricting **A** to be idempotent often allows us to find a polynomial time algorithm (Freese, Valeriote, Horowitz, K.).
- Big goal: Characterize, the complexity of all versions of the problem for all Maltsev conditions.
- Today's goal: Characterize, the complexity of deciding *k*-ary weak near unanimity for idempotent algebras when *k* is fixed.

• A *k*-ary weak near unanimity (*k*-WNU) is any idempotent operation that satisfies

$$t(y, x, \ldots, x) \approx t(x, y, \ldots, x) \approx \cdots \approx t(x, x, \ldots, y).$$

- A *k*-WNU is an example of a Taylor term.
- Given **A** idempotent, we can test in polynomial time whether **A** has a Taylor term (Bulatov).

• We say that **A** has *n*-local *k*-WNUs if for every $\overline{r}, \overline{s} \in A^n$ there exists a term *t* such that

$$t(s_i, r_i, \ldots, r_i) = t(r_i, s_i, r_i, \ldots, r_i) = \cdots = t(r_i, r_i, \ldots, s_i).$$
for all *i*.

• Translating to relations:

$$\operatorname{Sg}_{\mathbf{A}^{kn}}\left(\begin{pmatrix}\overline{\overline{s}}\\\overline{r}\\\vdots\\\overline{r}\end{pmatrix},\begin{pmatrix}\overline{r}\\\overline{\overline{s}}\\\vdots\\\overline{r}\end{pmatrix},\ldots,\begin{pmatrix}\overline{r}\\\overline{r}\\\vdots\\\overline{\overline{s}}\end{pmatrix}
ight)$$

has a member $(\overline{a}, \overline{a}, \dots, \overline{a})^T$ for some $\overline{a} \in A^n$.

Observation

If **A** has k-WNU then it has n-local k-WNUs for all n.

Lemma

If **A** is idempotent, has a Taylor term and has n-local k-WNUs then **A** also has (n + 1)-local k-WNUs.

Corollary

We only need to check for the existence of a Taylor term and for all 1-local *k*-WNUs. This is polynomial for *k* fixed.

$$R = \operatorname{Sg} \begin{pmatrix} \overline{s} & \overline{r} & \overline{r} & \overline{r} \\ \overline{r} & \overline{s} & \overline{r} & \overline{r} \\ \overline{r} & \overline{r} & \overline{s} & \overline{r} \\ \overline{r} & \overline{r} & \overline{r} & \overline{s} \\ d & c & c & c \\ c & d & c & c \\ c & c & d & c \\ c & c & c & d \end{pmatrix}$$

•

Applying an *n*-local 4-WNU, we have

$$(\overline{a},\overline{a},\overline{a},\overline{a},\overline{a},b_1,b_2,b_3,b_4)^T \in R$$

The last four entries can be permuted.

Lemma (Barto, Kozik, Niven)

If **G** is a Taylor algebra, $E \leq_s \mathbf{G}^2$ and (G, E) has a component of algebraic length 1, then E contains a loop.

In our case, let

$$G = \{ (x_1, x_2) \colon \exists y, z, (\overline{a}, \overline{a}, \overline{a}, \overline{a}, \overline{a}, x_1, x_2, y, z)^T \in R \}$$
$$E = \{ ((x_1, x_2), (x_2, x_3)) \colon \exists z, (\overline{a}, \overline{a}, \overline{a}, \overline{a}, x_1, x_2, x_3, z)^T \in R \}.$$

 \overline{G} has a Taylor term, E is subdirect, and (G, E) contains a 3-cycle and a 4-cycle in the same component (\Rightarrow algebraic length 1). The loop lemma gives us e, f such that $(\overline{a}, \overline{a}, \overline{a}, \overline{a}, e, e, e, f) \in R$.



Apply 4-WNU for e, f and obtain a tuple $(\overline{a}, \overline{a}, \overline{a}, \overline{a}, b, b, b, b)^T$.

- The "local to global" idea works for many linear Maltsev conditions in the idempotent case...
- ... but Dmitriy Zhuk has recently shown that "local to global" does not work for the minority operation:

 $t(x,x,y) \approx t(x,y,x) \approx t(y,x,x) \approx y.$

- Question: How hard is it to decide if A idempotent has a minority?Related problem: Given A that satisfies a strong Maltsev condition
 - M, produce operations of **A** that witness M.

Thank you for your attention.