

Quadratic quasigroups in public-key cryptography

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- ▶ In this talk I will show how this scheme works, then I will describe the quadratic quasigroups and loops.

- ▶ **private key:** $(\mathcal{L}_1, \mathcal{L}_2, \mathcal{P})$.
 - ▶ $\mathcal{L}_1, \mathcal{L}_2$ are automorphisms of the vector space \mathbb{F}_2^n (regular matrices)
 - ▶ \mathcal{P} is an invertible map $\mathbb{F}_2^n \rightarrow \mathbb{F}_2^n$ represented by quadratic Boolean polynomials.

General MQ-scheme

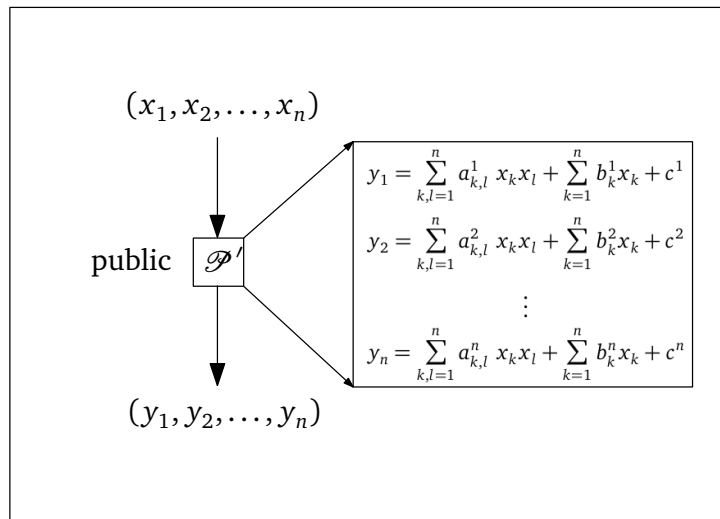
- ▶ **private key:** $(\mathcal{L}_1, \mathcal{L}_2, \mathcal{P})$.
 - ▶ $\mathcal{L}_1, \mathcal{L}_2$ are automorphisms of the vector space \mathbb{F}_2^n (regular matrices)
 - ▶ \mathcal{P} is an invertible map $\mathbb{F}_2^n \rightarrow \mathbb{F}_2^n$ represented by quadratic Boolean polynomials.
- ▶ **public key:** $\mathcal{P}' = \mathcal{L}_2 \circ \mathcal{P} \circ \mathcal{L}_1$.

Encryption and decryption

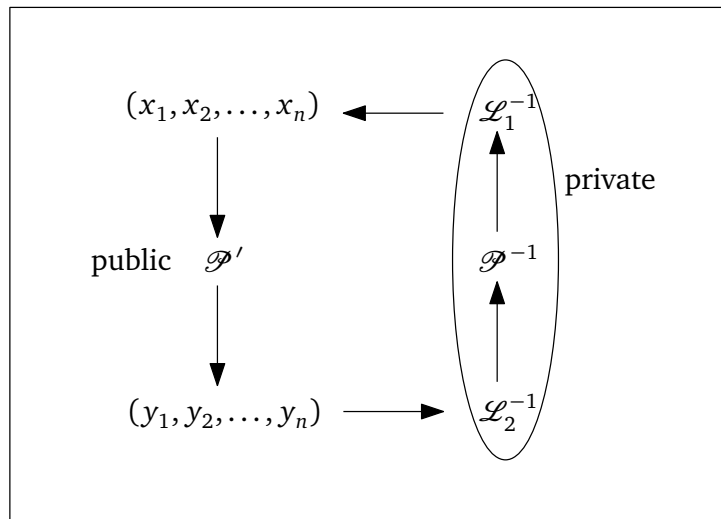
Secret message

$$(x_1, x_2, \dots, x_n) \in \mathbb{F}_2^n$$

Encryption and decryption



Encryption and decryption



Quadratic quasigroup

Let $(\mathbb{F}_2^m, *)$ be a quasigroup upon the vector space \mathbb{F}_2^m .

$$*: \mathbb{F}_2^{2m} \rightarrow \mathbb{F}_2^m, \text{ neco}$$

$$*: (x_1, \dots, x_m, y_1, \dots, y_m) \mapsto (z_1, \dots, z_m),$$

where $(x_1, \dots, x_m) * (y_1, \dots, y_m) = (z_1, \dots, z_m)$.

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$$*_i : \quad \mathbb{F}_2^{2m} \rightarrow \mathbb{F}_2$$

$$*_i : \quad (x_1, \dots, x_m, y_1, \dots, y_m) \mapsto z_i.$$

$*_i$ is a boolean map. There exists corresponding boolean polynomial f_i such that

$$\mathbf{x} * \mathbf{y} = (f_1(\mathbf{x}, \mathbf{y}), f_2(\mathbf{x}, \mathbf{y}), \dots, f_m(\mathbf{x}, \mathbf{y})),$$

for all $\mathbf{x}, \mathbf{y} \in \mathbb{F}_2^m$.

Quadratic quasigroup

Definiton

A quasigroup with an operation “ $*$ ” on the set \mathbb{F}_2^m is called *quadratic* if there exists quadratic polynomials $f_i \in \mathbb{F}_2[x_1, \dots, x_m, y_1, \dots, y_m], i = 1, \dots, m$ of the order at most 2 such that

$$\mathbf{x} * \mathbf{y} = \left(f_1(\mathbf{x}, \mathbf{y}), f_2(\mathbf{x}, \mathbf{y}), \dots, f_m(\mathbf{x}, \mathbf{y}) \right),$$

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for all $\mathbf{x}, \mathbf{y} \in \mathbb{F}_2^m$.

The map $\mathcal{P} : \mathbb{F}_2^n \rightarrow \mathbb{F}_2^n$ is constructed from quadratic quasigroups of order 2^5 in the following way:

Construction of the map \mathcal{P}

$$(x_1, \dots, x_5) = \mathbf{x}^1$$

$$(x_6, \dots, x_{10}) = \mathbf{x}^2$$

$$(x_{11}, \dots, x_{15}) = \mathbf{x}^3$$

$$(x_{n-4}, \dots, x_n) = \mathbf{x}^k$$

Construction of the map \mathcal{P}

$$(x_1, \dots, x_5) = \mathbf{x}^1 \rightarrow \mathbf{x}^1$$

$$(x_6, \dots, x_{10}) = \mathbf{x}^2 \rightarrow \mathbf{x}^1 *_1 \mathbf{x}^2$$

$$(x_{11}, \dots, x_{15}) = \mathbf{x}^3 \rightarrow \mathbf{x}^2 *_2 \mathbf{x}^3$$

$$\vdots$$

$$(x_{n-4}, \dots, x_n) = \mathbf{x}^k \rightarrow \mathbf{x}^{k-1} *_k \mathbf{x}^k$$

Construction of the map \mathcal{P}

$$\begin{array}{llll} (x_1, \dots, x_5) = \mathbf{x}^1 \rightarrow \mathbf{x}^1 & \rightarrow & \text{bijection } \mathcal{D} & \rightarrow \mathbf{y}^1 \\ (x_6, \dots, x_{10}) = \mathbf{x}^2 \rightarrow \mathbf{x}^1 *_{\mathbf{1}} \mathbf{x}^2 & \rightarrow & & \rightarrow \mathbf{y}^2 \\ (x_{11}, \dots, x_{15}) = \mathbf{x}^3 \rightarrow \mathbf{x}^2 *_{\mathbf{2}} \mathbf{x}^3 & \rightarrow & & \rightarrow \mathbf{y}^3 \\ \vdots & \vdots & & \vdots \\ (x_{n-4}, \dots, x_n) = \mathbf{x}^k \rightarrow \mathbf{x}^{k-1} *_{k-1} \mathbf{x}^k & \rightarrow & & \rightarrow \mathbf{y}^k \end{array}$$

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\vdots	\vdots		\vdots
$(x_{n-4}, \dots, x_n) = \mathbf{x}^k \rightarrow \mathbf{x}^{k-1} *_{k-1} \mathbf{x}^k \rightarrow$			$\rightarrow \mathbf{y}^k = (y_{n-4}, \dots, y_n)$

Inverting the map \mathcal{P}

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Properties of Gligoroski MQ-scheme

- Size of the keys (in KBytes):

n	Public key	Private key
140	168.7	9.8
160	251.6	11.3
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- Speed comparison (in 1000 cycles):

Algorithm	Sign	Verify
DSA 1024-bit	1041	1246
ECC 1024-bit	2147	4220
RSA 1024-bit	2939	99
MQQ 160-bit	10	140

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- ▶ Possible improvement/modifications:
 - ▶ Generate more complex quadratic quasigroups.
 - ▶ Use just left quadratic quasigroups.
 - ▶ Exclude some polynomials to make MQQ non-injective. Can be still used for signing.

Decomposition of quadratic quasigroup

Let $(\mathbb{F}_2^n, *)$ be a quadratic quasigroup and (f_1, \dots, f_n) its representation.

$$f_1 = \sum_{k,l=1}^n a_{k,l}^1 x_k x_l + \sum_{k,l=1}^n m_{k,l}^1 x_k y_l + \sum_{k,l=1}^n b_{k,l}^1 y_k y_l + c^1$$

$$f_2 = \sum_{k,l=1}^n a_{k,l}^2 x_k x_l + \sum_{k,l=1}^n m_{k,l}^2 x_k y_l + \sum_{k,l=1}^n b_{k,l}^2 y_k y_l + c^2$$

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$$\alpha(\mathbf{x}) \qquad \gamma(\mathbf{x}, \mathbf{y}) \qquad \beta(\mathbf{y}) \qquad \mathbf{c}$$

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$$\mathbf{x} * \mathbf{y} = \alpha(\mathbf{x}) + \gamma(\mathbf{x}, \mathbf{y}) + \beta(\mathbf{y}) + \mathbf{c}$$

Decomposition of quadratic quasigroup

- ▶ γ is a *bilinear map* $\mathbb{F}_2^n \times \mathbb{F}_2^n \rightarrow \mathbb{F}_2^n$ represented by bilinear forms.

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- ▶ Denote $\mathbf{o} = (0, 0, \dots, 0)$. Then

$$R_{\mathbf{o}}(\mathbf{x}) = \mathbf{x} * \mathbf{o} = \alpha(\mathbf{x}) + \mathbf{c}, \quad \text{a} \quad L_{\mathbf{o}}(\mathbf{x}) = \mathbf{o} * \mathbf{y} = \beta(\mathbf{y}) + \mathbf{c}.$$

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- ▶ α and β are bijective maps - *quadratic permutations*.

Decomposition of quadratic quasigroup

Theorem

For every quadratic quasigroup $(\mathbb{F}_2^n, *)$ there exist uniquely determined quadratic permutations α, β , a bilinear map γ and a vector \mathbf{c} , such that

$$\mathbf{x} * \mathbf{y} = \alpha(\mathbf{x}) + \gamma(\mathbf{x}, \mathbf{y}) + \beta(\mathbf{y}) + \mathbf{c},$$

for every $\mathbf{x}, \mathbf{y} \in \mathbb{F}_2^n$.

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$$\mathbf{x} * \mathbf{y} = \alpha(\mathbf{x}) + \gamma(\mathbf{x}, \mathbf{y}) + \beta(\mathbf{y}) + \mathbf{c}$$

Quadratic loop

- ▶ How do they look like?

Quadratic loop

Corollary

- ▶ Suppose that $(\mathbb{F}_2^n, *)$ is a quadratic loop with a unit \mathbf{o} .
- ▶ Then for every $\mathbf{x}, \mathbf{y} \in \mathbb{F}_2^n$

$$\mathbf{x} * \mathbf{y} = \mathbf{x} + \gamma(\mathbf{x}, \mathbf{y}) + \mathbf{y}.$$

- ▶ Such a loop is uniquely determined by the bilinear map γ .

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- ▶ Such a loop is uniquely determined by the bilinear map γ .
- ▶ It is easy to generate γ for a small n .

Quadratic loops of order 8

- ▶ I generated all quadratic loops of order 8.
- ▶ The number of these loops is 4384.
- ▶ 13 main classes.
- ▶ Contains all groups except the cyclic one.